Introduction to Differential Privacy

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Overview

- Need for Computational Privacy
- Plausible Deniability (Randomized Response)
- Differential Privacy
- Applications

Netflix Challenge

Anonymity is not enough!

In 2006, Netflix announced a \$1 Million prize challenge for the best <u>collaborative filtering</u> <u>algorithm</u> to predict user ratings. They released an anonymous version of their dataset.

In 2007, 2 researchers from UT Austin were able to de-anonymize the dataset using the open IMBD database.

Netflix Challenge



Need for Computational Privacy

Methods used:

- Distributed Computation
- Encrypted Computation
- Data swapping
- K-Anonymity
- Anonymization
- Rule Hiding

Need for Computational Privacy

Goal: Privacy-preserving Data Analysis Motivating Example: Census bureau

Adversary: Membership inference attack Data reconstruction attack Linkage attack

Intuition: Uncertainty in the process means uncertainty for the attacker

We need a mathematical guarantee on the "process" which helps us quantify and upper bound our loss of privacy

Need for Computational Privacy

Statistical analysis which learns that smoking causes cancer 2 levels of harms for each smoker:

1. Harm caused by smoking – what statistical analysis can help with

2. Harm caused by insurance companies becoming aware that person X is a smoker – higher insurance fee

We want to learn that "smoking causes cancer" to irradicate harm 1, without causing harm 2 to people in the process of data analysis.

Did you vote for the BJP?



Differential Privacy

Differential privacy is a system for publicly sharing information about a dataset which masks individual contributions while retaining the big picture, by adding some random noise to the data.

- Doesn't require attack modeling
- Privacy loss is quantifiable
- Compose multiple queries
- Accessible, minimal utility loss, easy to compute



Example



Example



Example



Basic Pipeline



Differentially Private Pipeline



Differential Privacy

A randomized algorithm M gives ε -differential privacy if for all pairs of data sets d, d' differing in the data of any one person, and all outputs S $Pr[M(d) = S] \leq e^{\varepsilon} Pr[M(d') = S]$

Where ε (+ve real number) is the controllable privacy budget parameter. The smaller its value, the better privacy guarantee you achieve.

Symmetric formulation

If a bad event is very unlikely when I'm not in the dataset (y) then it is still very unlikely when I am (x)

[2006, Cynthia Dwork, Frank McSherry, Kobbi Nissim and Adam D. Smith]

Randomized Response

With 50% probability --> BJP voters will say the truth & say yes With 50% probability --> BJP voters will give a random answer 25% -->Yes 25% --> No => BJP voters will say Yes with a 75% chance. P[M(BJP voter) = Yes] = 0.75P[M(BJP Voter) = No] = 0.25P[M(BJP non-voter) = Yes] = 0.25P[M(BJP non-voter) = No] = 0.75

Randomized Response

P[M(BJP voter) = Yes] = 0.75 | P[M(BJP Voter) = No] = 0.25P[M(BJP non-voter) = Yes] = 0.25 | P[M(BJP non-voter) = No] = 0.75 $Pr[M(d) = S] \le e^{F} Pr[M(d') = S]$

0.75 / 0.25 = 3=> $e^{\epsilon} = 3$ => $\epsilon = \ln(3) \sim 1.1$

Randomized Response

Randomized response offers a guarantee of $(\varepsilon = 1.1)$ - Differential Privacy.

This means that an adversary who thinks their target is in the dataset with probability 50% can increase their confidence to at most 75%.



Understanding privacy budget





Basics of DP



Laplacian Mechanism

Sample noise from the Laplace distribution and add that noise to your data.

Mean = o b = $\Delta f / \epsilon$

$$f(x \mid \mu, b) = rac{1}{2b} \expigg(-rac{|x-\mu|}{b}igg)$$



Laplacian Mechanism

No. Of votes without target = 1000

The adversary wants to know whether target user voted for BJP

Noisy result --> No. Of BJP votes with target = 1003

Blue curve --> True BJP vote count with target = 1001 Orange curve --> True BJP vote count with target = 1002

Blue curve is more likely than orange curve by a probability of e^ɛ



Query Composition

DP gives us the ability to compose multiple queries, with the privacy budgets linearly adding up, making it a weaker guarantee of privacy, but predictable.

If algorithm M1 is ε_1 -DP and algorithm M2 is ε_2 - DP, then publishing the result of both is ($\varepsilon_1 + \varepsilon_2$)-DP

Combined result C = (M1(d), M2(d)). This is because M1 and M2 are independent.

Properties

- No longer need an adversary model: You protect all info about an individual, & it doesn't matter what the adversary knows about you beforehand
- Privacy loss is quantifiable: greatest possible info gain
- Future-proof: robust to post-processing
- Automatically yields group privacy: kɛ for groups of size k
- Understand behavior under composition: Can bound cumulative privacy loss over multiple analyses
- Programmable Complex private analyses from simple private building blocks

Applications

- AI in healthcare -> sensitive patient information can help improve diagnosis of various diseases
- Usage statistics in Google Chrome using RAPPOR
- Contact tracing beyond encrypted bluetooth messages
- Model-centric Federated Learning for any ML based prediction
- Census Bureau
- IoT: Heartrate monitors
- Combating memorization in Neural Networks

References

- Gautham Kamath's course on Algorithms for Private Data Analysis: <u>http://www.gautamkamath.com/CS86o-fa2020.html</u>
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Thank

you :)